

INTERNATIONAL (MULTIPLE LANGUAGE/ NON- ENGLISH) DOMAIN NAME  
AND EMAIL USER ACCOUNT ID SERVICES SYSTEM

5 BACKGROUND OF THE INVENTION

This invention relates to multi-lingual domain names, translation between multi-lingual and normal Internet domain names and domain name routing protocols on the Internet. It also relates to Unicode (see The Unicode Standard, Version 3 by Unicode Consortium, The Unicode Consortium, Addison-Wesley Publishing Company, published  
10 February 11, 2000, 1072 pp.)

Throughout man's history there have been interconnecting networks, some of which include the situation of a language group of one network requiring translation to move smoothly onto another node or network. Examples of these include the road  
15 systems of the Inca Kingdom in the region of Peru, of the Roman Empire, and of China under various dynasties. More recently, by 1838, Samuel Morse introduced long - distance wired binary communication, Morse code and the telegraph. With the advent of the telephone and the computer, modern communication networks materialized.

Today the Internet is worldwide communication network. In a sense it originated  
20 in the United States and so has taken on a flavor of the English-speaking word in its naming protocols used to designated destinations and sources. However, the majority of people of the world use, or at least are fluent in, non-English languages. If one subtracts off the populations of the United States, the United Kingdom, Ireland, Canada, and Australia, one is still left with five-plus billion total world population minus three-tenths  
25 billion, or some 4.7-plus billion non-native English-speakers.

The Internet is fast becoming a significant market place in itself as well as providing a commercial information network. The Internet has opened up worldwide interconnectivity and information sources. Non-English speaking individuals and commercial firm of non-English speaking countries are significant players on the  
30 Internet. Most international corporations have web pages available on the Internet. It should be noted, however, that international companies tend to have different names, which consist of different characters, in different languages.

Recognizing that the Internet is still very much part of an English-speaking technical world, it is desirable to expand the domain name (DN) system to include multi-  
35 language (ML) domain names, while retaining as much as possible the substance and the flavor of the current Domain Name Specification as described in Network Working Group's Request for Comments (RFC) 1035 and RFC 1034. A reliable and flexible method for multi-language domain names (ML-DN) is needed. This method for the multi-language domain names should fit into the current system without producing name  
40 conflicts or requiring major restructuring of the Internet or its protocols. Ideally it should require a minimum new resources for its implementation.

Similar and parallel requirements apply as well to the multi-language e-mail user account identification (ML-UID) as to the multi-language domain name (ML-DN), where the regular e-mail account identification is ID. Also, such requirements apply to the  
45 multi-lingual URLs.

Kapoor (US Patent 5884038) indicates a method for providing an Internet protocol address with a domain name server based on a static array which apportions the web servers among the client domains that most frequently access the web servers. Nitta et al. (US Patent 4,641,264) shows a method of automatic translation between natural  
50 languages utilizing parsing and lexicon storage look-up. Kollin et al. (US Patent 4,774,665) invented a system for retrieving information from a plurality of remote database having at least two different languages. Information from a selected database is downloaded and then examined off-line. McGarvey (US 5,777,989) shows an enhanced name resolution for machines in several domains. The host is configured in more than  
55 one primary domain name server.

One approach to the multi-lingual Internet has been described by the iDNS working group (Asia Pacific Networking Group Internationalized Domain Name Server Working Group coordinating with the Internet Engineering Task Force). This approach is further described on the Internet at the site: "<http://www.iDNS.org/>". It translates the  
60 foreign non-English characters into Unicode ([www.unicode.org](http://www.unicode.org)) and then converts them to a double byte English equivalent ([www.idns.org](http://www.idns.org)). The iDNS working group itself has noted certain limitations in this approach. It may cause identical coding sometimes, i.e., an English domain name generated by the iDNA approach from a non-English charter

domain sometimes conflicts with an existing English based domain name. Also, the  
65 available length for the international language name is more limited since every  
international character in Unicode requires the equivalent of two English-number  
characters to accomplish the coding.

A more versatile naming system would be more compatible with the international  
growth of the Internet. The DNS (Domain Name System) is the on-line distributed  
70 database system used to map human-readable machine names into IP (Internet Protocol)  
addresses. DNS servers throughout the connected Internet implement a hierarchical  
namespace that allows sites freedom in assigning machine names and addresses. The  
DNS also supports separate mappings between mail destinations and IP addresses.

The approach presently used to resolve domain name is analyzed here. As shown  
75 in Figure 1a (prior art), the client workstation A 1001 wants to talk to the server  
workstation B 1002, but A only knows the domain name (DN) of B as  
“www.b-station.com” 1003. Consequently A 1001 needs to get B’s 1002 IP (Internet  
protocol) address 1004, first. As an example, it might be 111.222.333.444 1004. Next A  
1001 looks into its database and finds out that there is a DNS (Domain Name System)  
80 server D 1005, i.e., D’s IP address. Now A send a DN/IP (Domain Name to Internet  
Protocol) translation request (as network route 1-R 1006, 4-R 1007) to D 1005 with a  
request for the IP (Internet Protocol) address of the domain name DN “www.b-  
station.com” 1003. When D 1005 receives the request from A 1001, D 1005 then checks  
its DNS database. Before D 1001 looks up its database for an IP match, however, it  
85 checks to determine whether the DN is in a valid format as defined in RFC 1035, i.e., less  
than 255 characters long, only has lower case characters from the 26 English alphabet, or  
“.”, or “-”.

If the DN is in the correct format and if D 1005 finds the corresponding IP  
address, i.e., 111.222.333.444 1004, for the DN “ www.b-station.com” 1003, D 1005  
90 then returns that IP address to A 1001 (by network route 4-G 1008, 1-G 1009). Otherwise  
if D 1005 cannot match the domain name (DN) to an Internet Protocol (IP) address, D  
1005 will continue on and interrogate other DNSs 1011 (e.g., by network route 6 1010).  
If no match is found, an error message or error code is returned to A 1001. Once A 1001

has received the returned IP address (111.222.333.444 **1004**) of B **1002** from D **1005** it  
95 can start talking to B **1002** (by network route 1-B **1012**, 5-B **1013**).

An internet is a network which comprises at least the Internet.

## SUMMARY OF THE INVENTION

100 The aim of this invention is to place those billions of people on this planet who  
are not English-speaking but want to have, or access, an appropriate domain name, on an  
equal footing with the English speaking populations.

The invention proposed herein employs a translation database which has an entire  
set of international domain names, which can be in any suitable coding format, as an  
105 entry corresponding to a valid, by current domain name definition, English based domain  
name or an IP (Internet Protocol) address. This method, advantageously, fits into the  
current domain name system, or DNS, system without generating conflicts, such as  
producing an unwanted duplicated domain name (DN). This method also has flexibility,  
advantageously, to work with any language and character set.

110 The DN/ML-DN to IP resolution scheme forms the basis of the invention. The  
ML-DN client does not talk to normal DNS, except in the special case when the ML-DN  
is actually a DN(normal English one). The normal DNS does not understanding ML-DN  
and cannot translate a ML-DN request to a normal DN request. The ML-DN translates  
the ML-DN request to a normal DN request. Then either the ML-DNS server or ML-DN  
115 client can send the translated ML-DN, i.e. a normal DN to a normal DNS to get the  
corresponding IP address. The choice of server or client is based on the considerations of  
optimization of speed and ease of programming. Both approaches can be used.

A fully functional ML-DNS will recognize DN as well. When ML-DNS is widely  
accepted, DNS will then be a subset of ML-DNS.

120 Similarly, a multi-lingual email user account ID (ML-MUID) can also be used  
with the use of a user account translating table (database). The creation of new account  
directories is not needed.

Again, similarly, a multi-lingual URL (universal resource locator) can also be  
incorporated with the Internet system utilizing a resource locator translator so the Internet

125 system will have available to, and can usefully handle, both the regular URL's and the multi-lingual ML-URL's.

130 BRIEF DESCRIPTION OF THE DRAWINGS

The above and other features and advantages of the invention will be more apparent from the following detailed description wherein:

Figure 1a (prior art) illustrates the method of the DN system in the network  
135 environment;

Figure 1b shows a working method of the ML-DN system in the network environment;

Figure 2 shows a block diagram of the working method of the ML-DN System on the Client Side;

140 Figure 3 shows the block diagram of the working method of the ML-DN System on the Server Side;

Figure 4 shows the block diagram of the working method of the ML-Email UID System on the Client Side;

145 Figure 5 shows the block diagram of the working method of the ML-Email UID System on the Server Side;

Figure 6 shows the ML-DN/DN resolving scheme.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

150 The following description is the best mode presently contemplated for carrying out the invention. This description is not to be taken in a limiting sense, but is merely made for the purpose of describing the general principles of the invention. The scope of the invention should be determined with reference to the claims.

The multi-language domain name (ML-DN) system enables the usage of multiple  
155 language domain names interwoven seamlessly into the current system based on the

current domain name server system. The current domain names are defined in the Internet Request For Comment document, RFC 1035. The current domain name is English (or Latin) letter-based only name, which utilizes only alphabet letters a to z, numbers 0 to 9 and symbols "." and "-".

160 The description of the preferred embodiments of the multiple language (ML) domain name system (DNS) starts with Figure 1b. Figure 1b shows an embodiment of the invention for resolving both domain name system (DNS) and multiple language domain name system (ML-DNS) domain names.

165 The ML-DNS is indicated as server D' 1501 on Figure 1b. In this embodiment of the invention, B 1002 has an international domain name, for example, "BEB.BEΦ.KΛΓ.KΛ " 1502 as well as its English language domain name (DN): "www.b-station.com". If A 1001 wants to talk to B 1002 and A 1001 prefers to use the name "BEB.BEΦ.KΛΓ.KΛM" 1502 it cannot use the normal DNS, since it won't work with the invalid name "BEB.BEΦ.KΛΓ.KΛM " 1502. However, in this  
170 embodiment of the invention, A 1001 makes use of the ML-DNS (multi-lingual domain name server) D' 1501. The client workstation A 1001 sends the ML-DN/IP translation request to D' 1501 (by network route 1-Y 1503, 2-Y 1504). Upon receiving the request, D' 1501 first checks 3001 to determine if the request is a normal one, i.e., a normal English domain name (as shown in Figures 2 and 3). Upon examination, should the DN  
175 be a normal English one, D' 1501 sends a DN/IP translation request to D 1005 and upon receiving the IP address, 111.222.333.444 1004 D' 1501 sends it back to A 1001. In a case like this, D' 1501 works just like a DNS relay server.

When the received DN is a ML-DN, the server checks to determine whether it's available for direct translation of the ML-DN to an IP (111.222.333.444). In this  
180 circumstance, the ML-DNS and DNS databases are effectively integrated. In another circumstance, the direct translation is not available. This is more often situation. Consequently, the server checks its ML-DN to DN translation database for a valid DN match, i.e. translates "BEB.BEΦ.KΛΓ.KΛM" to "www.b-station.com". When the server does not find a valid DN match it interrogates other ML-DNS, as is indicated by network  
185 route 6'.

Then the DN "www.b-station.com" will be passed to D 1005 as a normal DN/IP translation request (by network route 3-R 1507 as shown in Figure 1b). Once the IP (111.222.333.444) is received at D' 1501 (by network route 3-G 1508 as shown in Fig. 1b), it is returned to A 1001 (by network route 2-G 1509, 1-G 1009 as in Fig. 1b).

190 The ML-DN database, in a preferred embodiment, requires a unique coding system for each language character set. Because there are sometimes more than one coding system for each language character set, a converting processing is necessary in those cases for the ML-DNS system. In one embodiment, such a procedure can be implemented on the server side with special message formatting rules. For example, users  
195 may be required to use a language-coding flag string, or a particular pre-specified coding system. In a more preferred embodiment, this converting procedure is implemented on the client side to speed up the ML-DN/IP translation.

The Unicode is a preferred candidate for the implementation of the unique language coding system. Although the current 16-bit Unicode ends to result in large files,  
200 it can, advantageously, resolve many conflicts caused by other coding systems. One can see that the embodiment of the ML-DN system is set up as a symmetrical analog to the existing DN.

As shown in Fig. 2, the input DN string 2001 is examined on the client side before being sent out. If 2002 the DN string 2001 is a normal (English) one then it will be either  
205 sent directly to a normal DNS or it will be encoded with a language (English) indicating word 2006 before being sent out to the ML-DNS as a ML-DN/IP translation request. If 2003 the DN string 2001 is a ML-DN, then a converting process 2003 will be employed in order to convert the string coding to a unique one for the detected language 2004 character set. When the language character set has been detected 2004, the DN input  
210 string is translated to a Unicode representation 2005. Then the converted ML-DN string is encoded with a language indicating word, i.e., a language flag header 2006, before being sent out to the ML-DNS 2007.

In an embodiment of this invention, a root ML-DNS database is maintained, worldwide. The root ML-DNS database is divided into sub-root database according to  
215 different language character set categories worldwide. The root ML-DNS databases hold all the registered ML-DNS, which are updated on a regular basis. All the other client ML-

DNS databases 1505 get the updates from the root databases (by way of route 6' 1506 in Fig.1b). The ML-DNS databases will work in a way analogous and parallel to the current DNS database.

220 The method for building and handling ML-DNS databases represent an enhancement and improvement of the existing system of DNS database methods (RFC 1035, RFC 1183 and RFC 1664). The invention comprises an improvement over the current DNS system. The ML-DNS root databases are not typically a single database residing at one physical site. The database usually consists of multiple sub-root databases  
225 maintained on different sites, even, for example, in different countries. In referring to a singular database, the usage, here, is meant to apply, also, to this plurality of data repositories. One embodiment of the invention has a format of the ML-DNS ML-DN/IP mapping set of databases as follows. For each ML-DNS sever there is a set of database for holding the multiple language Resource Records (ML-RR's as an enhancement of  
230 RR's for normal DNS). The types of the ML-RR's include as enhancements of ordinary DNS resource records SOA, NS, and AM. The SOA (Start Of Authority) is a keyword used with DNS to denote the beginning of the records for which a particular server is the authority. Other records in the server are reported as non- authoritative. SOA (Start Of Authority), NS (Name Server) records and AM (Address Mapping) records for the set of  
235 DNS databases have the counterparts: (1) ML-Start-of-Authority (ML-SOA) Records, (2) ML-Name Server (ML-NS) Records, (3) ML-Address Mapping (ML-AM) Records. The formats for the DNS RR's and zone files are defined in RFC 1035.

The syntax used in ML-DNS database files is an enhancement of that used in the DNS database, where the record fields for the ML-DN are be in a Unicode format, i.e., a  
240 16-bit hex address format (xxxx, where x=0-9, A-F) for every character in the ML-DN and where “.” is used as the separation symbol for the ML-DN labels instead of “.” as in normal DN's. In addition, the corresponding types for labeling the different records-SOA, NS and AM are replaced by ML-SOA, ML-NS and ML-AM. For the ML-DNS there is one master zone file, which have a ML-SOA record with the following format:

245  
<owner><ttl> IN ML-SOA<zone file source><mailbox>(<serial><refresh><retry>  
<expire><m-ttl>)



which is an exact duplicate of SOA for DNS except with the string ML-SOA instead of the string ML.

250

The ML-NS records in the following format, resembling NS records for the DNS with NS replaced by ML-NS:

```
<owner> <ttl> IN ML-NS<host.domain-name>
```

255

The control entries for designating the file names of ML-AM (multi-lingual address mapping) record databases for every valid language character set are:

260

```
$ORIGIN <domain name> [<comment>]
```

```
$LANGUAGE<language-character-set identifier><domain-name in Unicode format for  
current language-character-set> [<comment>]
```

265

```
$INCLUDE<file name> <language-character-set identifier>[<domain-  
name>][<comment>]
```

Then in the database file for a given language-character-set, the ML-AM (multi-lingual address mapping) records will have the following formats:

270

```
<domain-name in Unicode for current language-character-set>
```

```
<language-character-set identifier> IN ML-AM <valid English domain-name>
```

The ML-DNS database can be constructed following the above method examples using reliable and logical record formats to insure a reliable service.

275

The flow of the translation process for converting the "b-station's" Russian DN (in Unicode) to a normal IP address is described below.

280 Russian DN (Unicode)-> (ML-DNS)->normal English DN (non-Unicode) ->(DNS)->IP  
address

The process {Russian DN (Unicode)->(ML-DNS)->normal English DN} is  
discussed in detail in terms of a level-by-level, i.e., domain level, top-level-domain (tld)  
285 on down, resolving method. Each time one level of the ML\_DNS domain name is  
resolved, i.e., converted to a normal English DN fragment, the normal English DN  
replaces that ML-DNS DN fragment. In order to tell what part of the ML-DNS name  
string has already been resolved the “..” (“double dot”) fragment serves as a separator  
between the resolved part and the unresolved part. That unresolved part is further  
290 resolved. This can happen in the same or another ML-DNS server. A Top Level ML-  
Domain Name translation database (the size of which is relatively small) can be placed  
on every ML-DNS server in order to speed up the translation process.

The working method for the ML-DN resolver is an enhancement of the working  
method of a normal DN resolver. The ML-DN resolving process is as follows. The local  
295 ML-DNS resolver first checks the local ML-DNS database or cache files (after the initial  
string code converting process 2005, Fig.2) to determine if it is possible to map the entire  
ML-DN 3001 (Figure 3) to normal DN 3002 locally. See Figure 3. If not 2003, it then  
queries a specified ML-DNS server 3004. The ML-DNS system is configured such that  
every ML-DNS 3004 server is able to translate the Multi-Lingual-Top Level Domain  
300 (ML-TLD) names to the normal TLD's. Since the number of TLD's is quite limited the  
ML-DNS can be configured in this manner. In the case where there is an increase in the  
number of TLD's, one TLD ML-DNS can be set up for the ML-TLD/TLD mapping).

When a ML-DN listener receives a DNS request 2001, it first tries to determine if  
the request is a normal (English) request 3001. If it is, then a normal DNS 3002 server  
305 can translate the domain name to an Internet Protocol address 3003. If the request is not a  
normal (English) request 2003, then a ML-DNS server is queried as to its availability for  
direct translation of the ML-DN to an IP address. If it is available for this translation  
3007, then that translation is carried out 3008 and a real IP address will be returned 3003.  
If it is not available 3005 for direct ML-DN/IP translation, the first step is to translate

310 from ML-DN to DN **3011**. Once the translation for the DN **3011** is achieved, then the  
DN is sent **3006** to the DNS **3002** for translation from DN to IP, which is then sent on  
**3010** to return a real IP address.

There is a ML-DNS **3004** server under every TLD. For example, there is a  
“mldns.com” for the COM domain, which is responsible for mapping the second level  
315 part of the domain name. And for every second level domain there is another ML-DNS  
server **3008**, and so on. A physical server can host several domains. Thus, the number of  
servers needed for ML-DN to/from DN mapping is not necessarily large. The iterative  
process **3001**, **3004**, **3006**, **3002**, **3008**, **3003** will be carried on until the entire ML-DN  
has been mapped to the DN **3009**, **3010**. After the mapping, the normal DNS query will  
320 be made to get the desired IP address **3003**.

The IP packet format **3008** of ML-DNS queries can take the same form as for the  
DNS queries with the use of the appropriate type in the <Type> and <> fields.

An example of an embodiment of the ML-DN-/DN resolving method is shown  
on Figure 6. The example treats a case where the ML-DNS database maps mapping  
325 “BEB.BEΦKΛΓ.KAM” **6001** to “www.b-station.com” **6023** assuming the Unicode IP  
address: 111.222.333.444. The first “..” is part of the translated code while the second  
“..” may represent the separation between the first “translated” and the second “un-  
translated” part. In this example, the Russian language domain name (ML-DN) is  
“BEB.BEΦKΛΓ.KAM” and has the ML-DN designated address of  
330 “111122221111..111122223333444455556666..444477778888” **6001**.

In the boot zone file listings below, please note that “;” indicates the start of a  
comment. First the domain name to be resolved **6001** is checked to determine if it can be  
resolved locally **6002**. If it can **6003**, one has a DNS query on “www.b-station.com”  
**6023**. Next, a check is made to determine whether the primary ML-DNS **6005** can  
335 resolve the domain name. If it can **6006**, then one forwards the query to “www.b-  
station.com” **6023**.

If the primary ML-DNS **6005** cannot resolve the domain name, then one  
progresses on to TLD **6007** (top level domain) resolving. If it cannot **6008**, then an error  
**6013** is returned. If the primary ML-DNS **6005** can begin to resolve the ML-DN it  
340 proceeds to do so. The resolution of the last 12 integers “444477778888” to the top-level

domain (TLD) “com” is shown **6010**. The resolving continues on **6011**. If no further resolution can be done **6012**, then an error is returned **6013**. If further resolution can be done **6024**, the next resolution **6014** is that of “111122223333444455556666” to “b-station”. The resolution so far is:

345

“111122221111..111122223333444455556666..444477778888

to

350

“111122221111..b-station..com”.

Next, one determines if more resolving on ml-dns.com **6015** can be done. If yes, **6019**, one proceeds to resolve “111122221111.b-station.com” to “www.b-station.com” **6022**. If no further resolving **6016** can be done on “ml-dns.com”, then one attempts resolving “ml-dns.b-station.com” **6020**. If no further resolution **6017** can be done, an error is returned **6013**. If further resolution can be done, then one proceeds to further resolve “111122221111” into “www” **6022**. In all successful ML-DNS name resolutions, one then proceeds as a DNS query on “www.b-station.com” **6023**.

355

The iterative query chain starts on the DNS server for TLD “COM” i.e., the top level domain “.com”.

360

Boot zone file on mldns.com

Com            IN        ML-SOA mlns.com    root.mlns.com  
(169712        3200    7200    940000        96000); ML-SOA record

365

IN        ML-NS            mlns.com ; ML-NS record

\$ORIGIN com

\$LANGUAGE            110 444477778888

370

\$INCLUDE ru\_name.db 110; language identifier for Russian

“111122221111..b-station..com”

Then in file ru\_name.db one finds:

111122223333444455556666 110; IN ML-AM b-station

375

380

An inquiry is then sent to mldns.b-station.com:

Boot zone File on mldns.b-station.com

385

b-station.com IN ML-SOA mlns.b-station.com root-station.com  
(19712 1800 3600 640000 86000); ML-SOA record IN ML-NS

mlns.b-station.com ; ML-NS record

390

\$ORIGIN b-station.com

\$LANGUAGE 110 111122223333444455556666.444477778888

\$INCLUDE ru\_name.db 110; language identifier for Russian

395 then in file name ru\_name.db one finds:

111122221111 110; IN ML-AM www

Consequently, the name mapping has been accomplished.

400

## Email

The multiple language-enabling schemes can be applied to the Email system similarly. Figure 4 shows the client side multi-language user account ID (ML-UID) working scheme. The input string of email 4001 user ID is examined first 4002 before  
405 being used in the email header. If the ID is a normal one 4003 then it is used just as it is, or preferably in this embodiment, it can be encoded with a language (English) indicating flag word 4007 then added in to the mail header 4008. If the ID is a ML-UID 4004 then it will be converted to a unique coding format 4006 for the detected language 4005 character set in Unicode, encoded with a language indicating flag word 4007 and then  
410 added in to the mail header 4008.

As for the server side (as shown in Figure5) a ML-UID aware mail daemon 5001 is employed. The mail daemon 5001 checks the UID 5002 in the mail header in case the UID is a normal (English) one 5003 it just sends it for normal mail delivery 5006. If the UID turns out to be a ML-UID 5004, it checks a ML-UID to UID translation database  
415 5005 to find out the corresponding normal UID (in English), and then it sends the mail to the account with the found out UID 5006.

## URL

The Uniform Resource Locator (URL) is the combination of domain name (DN) sub domain name (sub-DN) and the server path name. For the ML-URL the ML-sub-DN  
420 will either converted to normal (English) sub-DN together with ML-DN or transferred to the server with the parent DN and then converted there, using the same converting mechanism as shown above. The conversion is the same by either route. The server path name part of the ML-URL (ML-URL path name) is treated the same way as for the ML-  
425 UID. The resulting string shall be passed to the server with the identified DN, where the converting of the ML-URL path name to normal URL path name by using a translating database. The server must handle multiple languages.

The database structures for mapping ML-UID ad path name part of ML-URL should be the same. As described above for ML-DNS Database. An enhancement uses  
430 two new types of records:

- ML-User ID (ML-UID) Records;
- ML-URL Path (ML-P) Records;

435           There is a master file for each of the databases, which includes the information for mapping from different language character set:

```
$Origin <domain name> [<comment>]
$LANGUAGE<language-character-set identifier><domain-name in Unicode format for
440 current language-character set> [<comment>]
$INCLUDE<file-name><language-character-set identifier>[<domain name>] [comment]
```

And the enhanced formats for ML-UID Records are:

```
445 <user ID in Unicode format for current language-character-set><language-character-set
identifier> IN ML-UID<Valid English user ID>
```

Then the formats for ML-P Records are:

```
450 <path name part of URL in Unicode format for current language-character-set>
<language-character-set identifier> IN ML-P<Valid English pathname part of URL.>
```

The process for resolving the ML path names will be localized on the server with the domain name. But the resolving scheme can also be iterative, as is the domain name resolving scheme. In other words the translation database can be placed under every

455 directory, which is designed to have sub-directories or files with ML names. Thus the full directory name can be translated level-by-level starting from the very top level to the very bottom level just by checking the database file holding the ML-P records for the items under the current level. The separation symbol “/” for separating different levels in the

460 pate name format can be replaced by “//” for a path name recorded only in 16-bit coding format for Unicode (shown in the text above). Also each database file can have multiple

entries mapping to the same normal path name string with the same language character set identifier.

For all the database mentioned above all the ML string are recorded in a 16-bit coding format (see above) by suggestion a more efficient method can be employed to reduce some unnecessarily long strings. Since, except for Chinese Japanese and Korean characters which all derive from ancient Chinese characters, most language character sets are “alphabetical” in nature, similar to English. Therefore the amount of necessary characters for these “alphabetical” character set is much limited, say, well under 256 different characters.

In that case the coding can be recorded as a base code (in 16-bit form) plus a shifting code (in 8-bit form). Then the string primarily can be recorded in 8-bit formats. Consequently the record size can be reduced greatly. Some control information may be included in the string for the implemented coding procedure.

There are many different coding systems for a number of languages, for example, Chinese has GB, Big5; Russian has KOI8-R, CP866; Japanese has SJIS, EUC; Korean has KSC5601; French has ISO-8859-1 and Finnish has ISO-8859-1. Consequently, in one embodiment of the invention, a system for converting different coding versions of a single language character set to a unique coding system is used. This can be done either on the server side or on the client side.

A modified version of UNIX Sendmail for ML-UID can be used to do the ML-UID/UID mapping. A modification for an APACHE Web Server can be programmed to enable the ML path mapping. The ML-UID database only resides on the appropriate Mail Server and ML-P-Web Server.

One method for simplifying the client part of the job that consists of determining the language character set is outlined here. There are certain characters in each language-character-set coding system, which or the combination of which, are distinctive for the same English string (in term of the meaning of the string). So we can use a specific ML-word as part of the ML-URL. This ML-word should appear the same, and mean the same thing, for the same language, but at the coding level it is different for different coding sets. The same coding system is to be used for the whole string of URL and/or the e-mail address. Thus, the language character set can be easily verified. For example, the URL



“www.b-station.com” in Russian may look like “BEB.BEΦKΛΓ.KΛM.”, the counterpart of the word “www” in Russian is defined as “BEB” which appears the same with any  
 495 Russian coding system. But at the real coding level it may looks “345678” for one Russian character set and “129876” for the other one. By using a small translation table on the client side it will be very easy to figure out what language character set is used by the user. The most foolproof method, of course, may be just to let the user identify what language character set he is using when he is entering the ML-URL, ML-UID or ML-DN.

500 All the above mentioned network activities are designed to be carried out in an IP network environment. But one should be able to use similar scheme with other network protocols. In terms of Server side programming the core technology may be based on UNIX BIND programs.

505 The multi-lingual method maps only the ML-DN to a valid normal DN. If one already has a normal (English) DN, the cost for setting up a ML-DN is much less. This name system can be run on any computer connected to the Internet. It will work independently or co-operatively with the current DNS for converting domain names to IP addresses.

510 While the invention herein disclosed has been described by means of specific embodiments and applications thereof, numerous modifications and variations could be made thereto by those skilled in the art without departing from the scope of the invention set forth in the claims.